

Predikatuen Logikako teorema batzuk

$\models \forall x (\phi \vee \neg \phi)$	<i>Hirugarren Baztertuaren Legea</i>
$\models \forall x \neg(\phi \wedge \neg \phi)$	<i>Kontraesan ezaren Legea</i>
$\models \forall x (\phi \rightarrow \phi)$	
$\models \phi [x/c] \rightarrow \exists x \phi$	<i>Orokortze existentziala</i>
$\models \forall x \phi \rightarrow \exists x \phi$	<i>Partikularizazio existentziala</i>
$\models \forall x \phi \rightarrow \phi [x/c]$	<i>Partikularizazioa</i>
$\models \forall x (\phi \rightarrow \psi) \rightarrow (\forall x \phi \rightarrow \forall x \psi)$	<i>Unibertsalaren distribuzioa baldintzan</i>
$\models \forall x (\phi \rightarrow \psi) \rightarrow (\exists x \phi \rightarrow \exists x \psi)$	
$\models \exists x \forall y \phi \rightarrow \forall y \exists x \phi$	

Predikatuen Logikako baliokidetzak batzuk

$\forall x \phi \equiv \neg \exists x \neg \phi$		$\forall x \phi \equiv \forall y \phi [x/y]$	Aldagaien
$\neg \forall x \phi \equiv \exists x \neg \phi$	Zenbatzaileen	$\exists x \phi \equiv \exists y \phi [x/y]$	aldaketa
$\forall x \neg \phi \equiv \neg \exists x \phi$	Interdefinizioak		
$\neg \forall x \neg \phi \equiv \exists x \phi$			

(y-k ez badauka ϕ -n agerpen askerik)

Distribuzioa:

$$\forall x (\phi \wedge \psi) \equiv \forall x \phi \wedge \forall x \psi$$

$$\exists x (\phi \vee \psi) \equiv \exists x \phi \vee \exists x \psi$$

Kuantifikazio hutsala

$$\phi \equiv \forall x \phi$$

$$\phi \equiv \exists x \phi$$

(x-ek ez badauka ϕ -n agerpen askerik)

Zenbatzaileen esportazioa baldintzan

$$\forall x (\phi \vee \psi) \equiv \forall x \phi \vee \psi$$

$$\exists x (\phi \wedge \psi) \equiv \exists x \phi \wedge \psi$$

$$\forall x (\psi \rightarrow \phi) \equiv \psi \rightarrow \forall x \phi$$

$$\exists x (\psi \rightarrow \phi) \equiv \psi \rightarrow \exists x \phi$$

$$\forall x (\phi \rightarrow \psi) \equiv \exists x \phi \rightarrow \psi$$

$$\exists x (\phi \rightarrow \psi) \equiv \forall x \phi \rightarrow \psi$$

(x-ek ez badauka ψ -n agerpen askerik)

(x-ek ez badauka ψ -n agerpen askerik)

Zenbatzaileen independentzia

$$\forall x \forall y \phi \equiv \forall y \forall x \phi$$

$$\exists x \exists y \phi \equiv \exists y \exists x \phi$$

$$\exists x \exists y A^2xy \equiv \exists x \exists y A^2yx$$

$$\exists x \exists y (A^1x \wedge A^1y) \equiv \exists x A^1x \wedge \exists y A^1y$$